

Automatic reformulation using optimization bridges

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Motivation

MathOptInterface (MOI)

MOI in a nutshell:

- `add_variable(model)`.
- `add_constraint(model, func, set)`, e.g. $2x + 3y = 1 \rightarrow (2*x + 3*y)\text{-in-EqualTo}(1.0)$.
- `set`, `get` attributes, e.g., `ObjectiveSense`, `ObjectiveFunction`.

Extensible framework:

- **Generic** on attribute, function and set types. New ones can be defined **independently**.
- Solver-**specific** features easily exposed to JuMP/MOI users through **custom** attributes.
- Expose **specialized** problem structure easily through **custom** functions, sets (e.g. Sum-of-Squares variables/constraints).

Semidefinite programming

$$\begin{array}{ll} \underset{Q \in \mathcal{S}^n}{\text{minimize}} & \langle C, Q \rangle \\ \text{subject to} & \langle A_i, Q \rangle = b_i \\ & Q \succeq 0 \end{array} \qquad \begin{array}{ll} \underset{y \in \mathbb{R}^n}{\text{maximize}} & \langle b, y \rangle \\ \text{subject to} & \sum_i A_i y_i \preceq C \end{array}$$

File format: SDPA

Solvers: CSDP, SDPA, DSDP, SDPLR, ...

Variables: Q block diagonal, nonnegative scalar variables (1×1 blocks) or SDP matrices.

Constraints: Affine equations.

Conic Modelling

```
using JuMP
model = Model(...)
@variable(model, -1 <= x <= 1)
@variable(model, y)
@variable(model, z <= 0)
@constraint(model, [x + y x
                    y      x - y] in PSDCone())
@constraint(model, [x + y, z, y] in SecondOrderCone())
@objective(model, x^2 - 2x*z + z^2)
```

The gap between models and solvers

The solver interface should only support structures and the algorithm **exploits**:

- n solvers and m structures $\rightarrow mn$ transformations $\rightarrow$ **unscalable** for large m, n .
- enables **evaluation** of formulation **quality**, e.g. automatic transformation and automatic dualization.

The model should

- be **independent** from solvers.
- represent the structure **exploitable** by algorithms.
- allow representable structure **unknown** to solvers, e.g. Sum-of-Squares variables/constraints.

Bridging the gap

$$\begin{aligned}x \in S_1 &\Leftrightarrow Ax \in S_2 & AS_1 &= S_2 \\ A^*y \in S_1^* &\Leftrightarrow y \in S_2^* & S_1^* &= A^*S_2^*\end{aligned}$$

In Lagrangian:

$$\langle Ax, y \rangle_2 = \langle x, A^*y \rangle_1$$

Transformation of variable-in- S_2 to variable-in- S_1 .

Primal Transform value v to Av .

Dual Transform dual y to $A^{-*}y$.

Transformation of f -in- S_1 constraint to Af -in- S_2 constraint.

Primal Transform value v of Af to $A^{-1}v$ of f .

Dual Transform dual y of A^*y .

Examples

FlipSignBridge

- Variable $x \geq l$ substituted by $x = -y$ where $y \leq -l$.
- Constraint $a^\top x \leq \beta$ transformed into $-a^\top x \geq -\beta$.

VectorizeBridge

- Variable $x \geq l$ substituted by $x = y + l$ where $y \in \mathbb{R}_+^1$.
- Constraint $a^\top x \leq \beta$ transformed into $[a^\top x - \beta] \in \mathbb{R}_-^1$.

IntervalBridge

- Constraint $l \leq a^\top x \leq u$ transformed into $a^\top x \geq l$ and $a^\top x \leq u$.

IndicatorBridge

- **Constraint** y binary, $\neg y \Rightarrow a^\top x \leq \beta$ transformed into z binary, $z + y = 1$ and $z \Rightarrow a^\top x \leq \beta$.

FreeBridge

- **Variable** $x \in \mathbb{R}$ substituted by $x = y + z$ where $y \in \mathbb{R}_+$ and $z \in \mathbb{R}_-$.

SlackBridge

- **Constraint** $f \in S$ transformed into $f = x$ for variable $x \in S$.
- **Objective** $\min f$ transformed into $\min x$ with $f \leq x$.

RSOC to PSD

Transforms rotated second order cone into positive semidefinite cone. By Schur Complement: $\|x\|_2^2 < 2tu, t, u > 0$ iff

$$\begin{cases} u > 0 \\ 2tu > x^\top x \end{cases} \Leftrightarrow \begin{cases} ul > 0 \\ t - x^\top (2ul)^{-1} x > 0 \end{cases} \Leftrightarrow \begin{pmatrix} t & x^\top \\ x & 2ul \end{pmatrix} \succ 0.$$

Dual of u

Diagonal at $(2,2), (3,3), \dots, (n,n)$ is $2u\mathbf{1}$ so dual of u is scalar product between $(2\mathbf{1})^\top$ and dual of diagonal at $(2,2), (3,3), \dots, (n,n)$.

Dual of x

Twice dual of $(1,2), (1,3), \dots, (1,n)$. Pay attention to scalar product of vectorized PSD cone.

Quadratic to RSOC

Transforms quadratic constraints into rotated second order cone constraints.

$$\frac{1}{2}x^T Qx + a^T x + \beta \leq 0$$

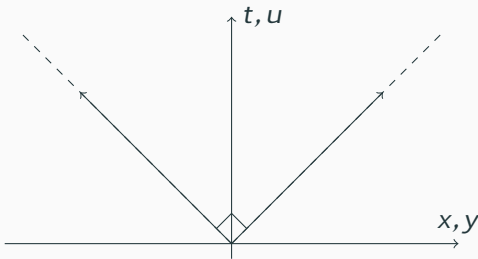
Q PSD $\rightarrow$ Cholesky $Q = U^T U$:

$$\|Ux\|_2^2 \leq 2(-a^T x - \beta)$$

so $(1, -a^T x - \beta, Ux) \in \text{RSOC}$.

Second order cone and complementarity slackness

If $(t,x), (u,y) \in \text{SOC}$ and $(t,x) \cdot (u,y) = 0$, then (u,y) is parallel to $(t,-x)$ or (t,x) is zero or (u,y) is zero.



Lemma

If $(1, s, x), (v, u, y) \in \text{RSOC}$ and $(1, s, x) \cdot (v, u, y) = 0$, then $y = -ux$ and $v = u \|x\|_2^2 / 2$.

Proof (sketch)

$$\begin{aligned} 0 &= (1, s, x) \cdot (v, u, y) \\ &= v + su + x \cdot y \\ (\text{Cauchy-Schwarz}) &\geq v + su - \|x\|_2 \|y\|_2 \\ (\text{RotatedSOC}) &\geq v + su - 2\sqrt{su} \\ (\text{AM-GM}) &\geq v + su - 2\frac{v + su}{2} \\ &= 0. \end{aligned}$$

See MOI QuadtoSOCBridge for details.

Quadratic to RSOC duals

Consider $z^T U^T U z / 2 + a^T z + \beta \leq 0$ transformed, so
 $(1, -a^T z - \beta, U z) = (1, s, x) \in \text{RSOC}.$

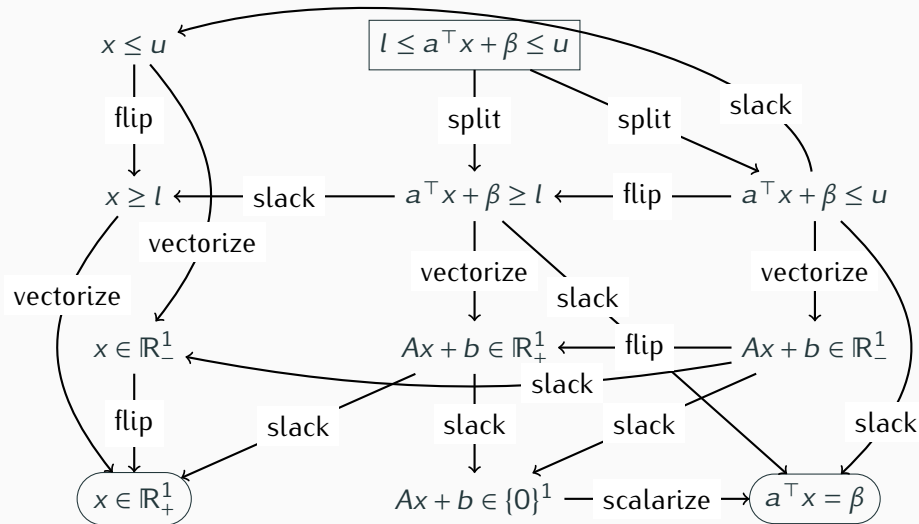
RSOC in Lagrangian:

$$\begin{aligned}(1, s, x) \cdot (v, u, y) &= u(1, s, x) \cdot (\|x\|_2^2 / 2, 1, -x) \\&= u(\|x\|_2^2 / 2 + s - \|x\|_2^2) \\&= u(-\|x\|_2^2 / 2 + s) \\&= u(-\|Uz\|_2^2 / 2 - a^T z - \beta) \\&= -u(z^T U^T U z / 2 + a^T z + \beta)\end{aligned}$$

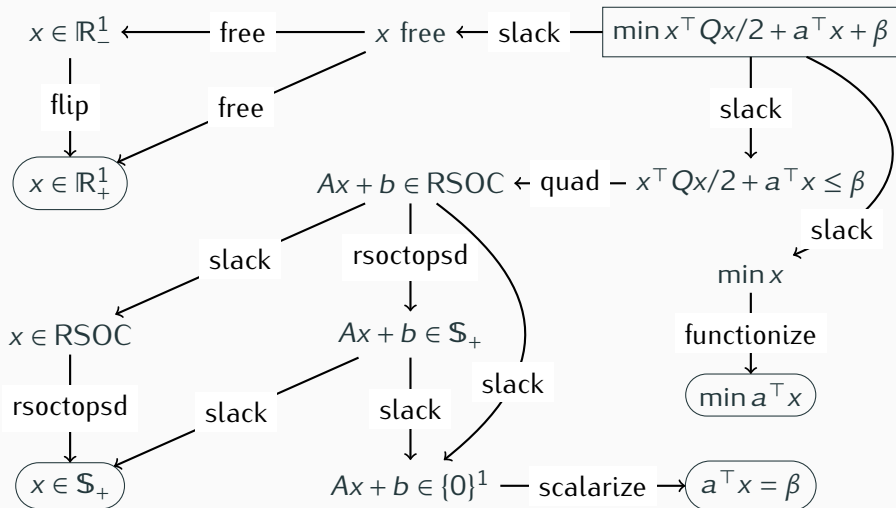
So the dual of the quadratic constraint is $-u$.

Combining bridges

Interval constraint for SDPA



Quadratic objective for SDPA



Automatic selection of bridges

Shortest path in directed Hypergraph

Nodes

Node for each set S (variable-in- S).

Node for each constraint F -in- S .

Node for each objective function F .

Types F and S are **not limited** to those defined in MOI.

Infinitely many nodes, we need to be **lazy**.

Edges

Each bridge defined possible **infinitely** many edges.

For each edge and ingoing node: outgoing nodes are

- variable-in- S created.
- constraints F -in- S created.
- objective F set.

Shortest Path Problem

Need to solve

$$d_V(S) = 1 + \min_{B \in \mathcal{B}_V(S)} \sum_{S' \in \mathcal{A}_V(B, S)} d_V(S') + \sum_{(F', S') \in \mathcal{A}_C(B, S)} d_C(F', S')$$

$$d_C(F, S) = 1 + \min_{B \in \mathcal{B}_C(F, S)} \sum_{S' \in \mathcal{A}_V(B, F, S)} d_V(S') + \sum_{(F', S') \in \mathcal{A}_C(B, F, S)} d_C(F', S')$$

$$d_O(F) = 1 + \min_{B \in \mathcal{B}_O(F)} \sum_{S' \in \mathcal{A}_V(B, F)} d_V(S') + \sum_{(F', S') \in \mathcal{A}_C(B, F)} d_C(F', S') \\ + d_O(\mathcal{A}_O(B, F))$$

with

- $d_V(S) = 0$ if constrained variables in S supported by optimizer.
- $d_C(F, S) = 0$ if F -in- S constraints supported by optimizer.
- $d_O(F) = 0$ if objective function F supported by the optimizer.

Shortest path algorithms ?

- Breath-First Search : For edges with cost 1
- Dijkstra : For edges with nonnegative cost
- Bellman-Ford : For edges with any real cost (+ negative cycles)

Choice: a modified *Bellman-Ford algorithm*.

Classical Bellman-Ford algorithm

- N: set of nodes
- E: set of edges
- d: distance
- b: next node

```
for _ in 1:length(N)-1:
    for each edge u=>v with weight w in E
        if d[u] + w < d[v]:
            d[v] = d[u] + w
            b[v] = u
        end
    end
end
```

Complexity $\mathcal{O}(|N| \cdot |E|)$

Target constraint types

Invariant: d_V , d_C , d_O are correct when defined.

constrained variable in S added by user

→ generate $\mathcal{V}$ (resp. $\mathcal{C}$): list of needed new entries to compute in d_V (resp. d_C).

F -in- S constraint added by user

→ generate $\mathcal{V}$ (resp. $\mathcal{C}$): list of needed new entries to compute in d_V (resp. d_C).

Objective function F constraint added by user

→ generate $\mathcal{V}$ (resp. $\mathcal{C}$, $\mathcal{O}$): list of needed new entries to compute in d_V (resp. d_C , d_O).

Add target variable type

Algorithm 1 recursive add S

```
add  $S$  to  $\mathcal{V}$ 
for  $B \in \mathcal{B}_{\mathcal{V}}(S)$  do
  for  $S' \in \mathcal{A}_{\mathcal{V}}(B, S)$  do
    if  $d_{\mathcal{V}}(S')$  not defined then
      recursive add  $S'$ 
    end if
  end for
  for  $(F', S') \in \mathcal{A}_{\mathcal{C}}(B, S)$  do
    if  $d_{\mathcal{C}}(F', S')$  not defined then
      recursive add  $F'$ -in- $S'$ 
    end if
  end for
end for
end for
```

Add target constraint type

Algorithm 2 recursive add F -in- S

```
add  $F$ -in- $S$  to  $\mathcal{C}$ 
for  $B \in \mathcal{B}_{\mathcal{C}}(F, S)$  do
  for  $S' \in \mathcal{A}_{\mathcal{V}}(B, F, S)$  do
    if  $d_{\mathcal{V}}(S')$  not defined then
      recursive add  $S'$ 
    end if
  end for
  for  $(F', S') \in \mathcal{A}_{\mathcal{C}}(B, F, S)$  do
    if  $d_{\mathcal{C}}(F', S')$  not defined then
      recursive add  $F'$ -in- $S'$ 
    end if
  end for
end for
```

Modified Bellman-Ford algorithm

```
while  $d_V$  or  $d_C$  changed do
  for  $S \in \mathcal{V}$ ,  $B \in \mathcal{B}_V(S)$  do
     $u \leftarrow 1 + \sum_{S' \in \mathcal{A}_V(B, S)} d_V(S') + \sum_{(F', S') \in \mathcal{A}_C(B, S)} d_C(F', S')$ 
    if  $u < d_V(S)$  then
       $d_V(S) \leftarrow u$ 
    end if
  end for
end for
for  $F\text{-in-}S \in \mathcal{C}$ ,  $B \in \mathcal{B}_C(F, S)$  do
   $u \leftarrow 1 + \sum_{S' \in \mathcal{A}_V(B, F, S)} d_V(S') + \sum_{(F', S') \in \mathcal{A}_C(B, F, S)} d_C(F', S')$ 
  if  $u < d_C(F, S)$  then
     $d_C(F, S) \leftarrow u$ 
  end if
end for
end while
```

Integration with JuMP

Extending JuMP macros

```
@constraint(model, [x + 1, x - y] in MOI.Zeros())
```

Implementation:

```
function build_constraint(  
    _error::Function,  
    func::Vector{<:AbstractJuMPScalar},  
    set::MOI.AbstractVectorSet)  
    return VectorConstraint(x, set)  
end
```

Extending JuMP macros: Custom set

```
@constraint(model, [x + 1, x - y] in SecondOrderCone())
```

Implementation:

```
function build_constraint(_error::Function,  
                        f::AbstractVector,  
                        s::AbstractVectorSet)  
    set = moi_set(s, length(f))  
    return build_constraint(_error, f, set)  
end  
function moi_set(::SecondOrderCone, dim::Int)  
    return MOI.SecondOrderCone(dim)  
end
```

Extending JuMP macros: PSD cone

```
using LinearAlgebra # For Symmetric
@constraint(model, Symmetric([x + 1 x - y
                             x - y y]) in PSDCone())
```

Implementation:

```
function build_constraint(_error::Function,
                        Q::Symmetric,
                        ::PSDCone)
    n = LinearAlgebra.checksquare(Q)
    func = [Q[i, j] for j in 1:n for i in 1:j]
    set = MOI.PositiveSemidefiniteConeTriangle(n)
    VectorConstraint(func, set,
                    SymmetricMatrixShape(n))
end
```

Sum-of-Squares example

Nonnegative quadratic forms into sum of squares

$$(x_1, x_2, x_3) \quad p(x) = x^T Q x \quad \text{unique}$$

$$x_1^2 + 2x_1x_2 + 5x_2^2 + 4x_2x_3 + x_3^2 = x^T \begin{pmatrix} 1 & 1 & 0 \\ 1 & 5 & 2 \\ 0 & 2 & 1 \end{pmatrix} x$$

$$p(x) \geq 0 \quad \forall x \iff Q \geq 0 \quad \downarrow \text{cholesky}$$

$$(x_1 + x_2)^2 + (2x_2 + x_3)^2 \longleftarrow x^T \begin{pmatrix} 1 & 1 & 0 \\ 0 & 2 & 1 \end{pmatrix}^T \begin{pmatrix} 1 & 1 & 0 \\ 0 & 2 & 1 \end{pmatrix} x$$

Nonnegative polynomial into sum of squares

$$(x_1, x_2, x_3) \quad (x_1, x_1 x_2, x_2) \quad \text{not unique}$$

$$p(x) = X^T Q X$$

$$x_1^2 + 2x_1^2 x_2 + 5x_1^2 x_2^2 + 4x_1 x_2^2 + x_2^2 = X^T \begin{pmatrix} 1 & 1 & 0 \\ 1 & 5 & 2 \\ 0 & 2 & 1 \end{pmatrix} X$$

$$p(x) \geq 0 \quad \forall x \iff Q \geq 0 \quad \downarrow \text{cholesky}$$

$$(x_1 + x_1 x_2)^2 + (2x_1 x_2 + x_2)^2 \longleftarrow X^T \begin{pmatrix} 1 & 1 & 0 \\ 0 & 2 & 1 \end{pmatrix}^T \begin{pmatrix} 1 & 1 & 0 \\ 0 & 2 & 1 \end{pmatrix} X$$

When is nonnegativity equivalent to sum of squares ?

Determining whether a polynomial is nonnegative is **NP-hard**.

Hilbert 1888

Nonnegativity of $p(x)$ of n variables and degree $2d$ is equivalent to sum of squares in the following three cases:

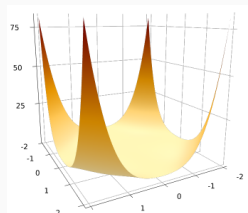
- $n = 1$: Univariate polynomials
- $2d = 2$: Quadratic polynomials
- $n = 2, 2d = 4$: Bivariate quartics

Motzkin 1967

First **explicit** example:

$$x_1^4 x_2^2 + x_1^2 x_2^4 + 1 - 3x_1^2 x_2^2 \geq 0 \quad \forall x$$

but is **not** a sum of squares.



Sum-of-Squares bridges

Polynomial $p \in \Sigma$ (p is SOS) iff $p = X^T Q x$ with $Q \in \mathbb{S}_+$ (Q is PSD). Hence $\Sigma = A\mathbb{S}_+$.

SOSPolynomialBridge: Transformation of variable-in- Σ to variable-in- $\mathbb{S}_+$.

Transformation of constraint F -in- Σ : SlackBridge + SOSPolynomialBridge.

Constraint Attribute

Examples: ConstraintPrimal, ConstraintDual, ConstraintFunction, ConstraintSet, ...

Redirected to bridge when constraint is bridged.

New attributes:

- GramMatrixAttribute: Gram matrix Q indexed by X .
- MomentMatrixAttribute: Moment matrix index by X , dual of constraint $Q \in \mathbb{S}_+$.
- MomentsAttribute: Vector of moments, dual of constraint $p = X^\top QX$.

Sum-of-Squares constraint macro

```
@constraint(model, p in SOS Cone())
```

Implementation:

```
function JuMP.build_constraint(_error::Function, p,  
                              cone::SOS Cone; kws...)
    coefs = coefficients(p)
    monos = monomials(p)
    set = JuMP.moi_set(cone, monos; kws...)
    shape = PolyJuMP.PolynomialShape(monos)
    return PolyJuMP.bridgeable(  
        JuMP.VectorConstraint(coefs, set, shape),  
        JuMP.moi_function_type(typeof(coefs)),  
        typeof(set)  
    )
end
```

Reshaping

Reshaping sets

```
function reshape_set(set::MOI.AbstractScalarSet,  
                    ::ScalarShape)  
    return set  
end  
function reshape_set(  
    ::MOI.PositiveSemidefiniteConeTriangle,  
    ::SymmetricMatrixShape  
)  
    return PSDCone()  
end
```

Reshaping polynomial sets

```
function JuMP.reshape_set(set::SOSPolynomialSet,  
                           ::PolyJuMP.PolynomialShape)  
    return set.cone  
end
```

Reshaping results

```
function reshape_vector(vectorized_form::Vector{T},
    shape::SymmetricMatrixShape) where T
    matrix = Matrix{T}(undef, shape.side_dimension,
        shape.side_dimension)

    k = 0
    for j in 1:shape.side_dimension
        for i in 1:j
            k += 1
            matrix[j, i] = matrix[i, j] =
                vectorized_form[k]
        end
    end
    return Symmetric(matrix)
end
```

Reshaping polynomial results

```
function JuMP.reshape_set(set::SOSPolynomialSet,  
                           ::PolyJuMP.PolynomialShape)  
    return set.cone  
end  
function JuMP.reshape_vector(x::Vector,  
                              shape::PolynomialShape)  
    return polynomial(x, shape.monomials)  
end  
function JuMP.reshape_vector(x::Vector,  
                              shape::MomentsShape)  
    return measure(x, shape.monomials)  
end
```

Dual shape

In JuMP:

```
dual_shape(shape::AbstractShape) = shape
```

In PolyJuMP:

```
function JuMP.dual_shape(shape::PolynomialShape)
    return MomentsShape(shape.monomials)
end
function JuMP.dual_shape(shape::MomentsShape)
    return PolynomialShape(shape.monomials)
end
```

JuMP and MathOptInterface provide an optimization framework **extensible** by **design**.

Bridges allow to communicate **custom** structure to solvers from **solver-independent** models.